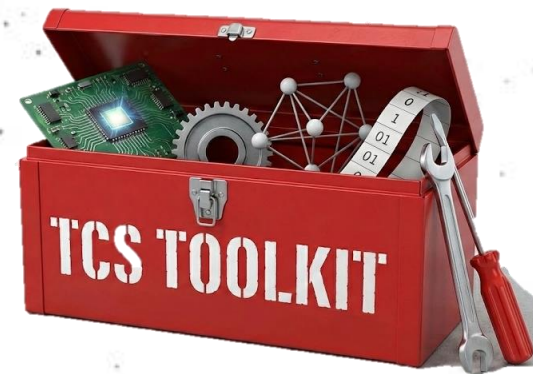


# CS 58500 – Theoretical Computer Science Toolkit

Lecture 19 (04/21)

Boolean Function Analysis (IV)

[https://ruizhezhang.com/course\\_spring\\_2026.html](https://ruizhezhang.com/course_spring_2026.html)



# Today's Lecture

- Influence and Isoperimetric Inequalities for the Hypercube
- KKL Theorem
- Unifying Talagrand and KKL

# Recall the Tensorization of Variance

**Theorem (Efron-Stein).** Suppose  $X_1, \dots, X_n$  are independent random variables. Let  $Z = f(X_1, \dots, X_n)$ . Then

$$\text{Var}[Z] \leq \mathbb{E} \left[ \sum_{i=1}^n \text{Var}_i[Z] \right]$$

where  $\text{Var}_i[Z](x_1, \dots, x_{i-1}, x_{i+1}, \dots, x_n) := \text{Var}_{X_i}[f(x_1, \dots, x_{i-1}, X_i, x_{i+1}, \dots, x_n)]$

In Boolean function analysis, we define the **influence** of  $f$  as:

$$I_i(f) := \mathbb{E}_{x_{-i}}[\text{Var}_i[f](x)], \quad I_f := \sum_{i \in [n]} I_i(f)$$

- If we define  $\partial_i f := f - \mathbb{E}_i[f]$ , then  $I_i(f) = \mathbb{E}_{x_{-i}} \left[ \mathbb{E}_i [f - \mathbb{E}_i[f]]^2 \right] = \|\partial_i f\|_2^2$

# Influence

For  $f: \{0,1\}^n \rightarrow \{0,1\}$ ,

$$\partial_i f(x) = f(x) - \frac{1}{2}(f(x) + f(x + e_i)) = \frac{1}{2}(f(x) - f(x + e_i))$$

$$I_i(f) = \frac{1}{4} \Pr[f(x) \neq f(x + e_i)]$$

## Example:

- Parity:  $I_i = \frac{1}{4}$  and  $I_{\text{Parity}} = \frac{n}{4}$

$$I_f = \frac{1}{4} \mathbb{E}[s_f(x)]$$

# Influence

For a general  $f: \{0,1\}^n \rightarrow \mathbb{R}$ ,

$$\begin{aligned} I_i(f) &= \frac{1}{4} \mathbb{E} \left[ (f(x) - f(x + e_i))^2 \right] \\ &= \frac{1}{4} \mathbb{E} \left[ \left( 2 \sum_{S \ni i} \hat{f}(S) \chi_S(x) \right)^2 \right] \\ &= \sum_{S \ni i} \hat{f}(S)^2 \end{aligned}$$

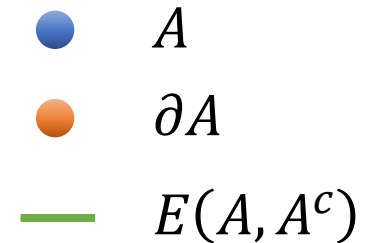
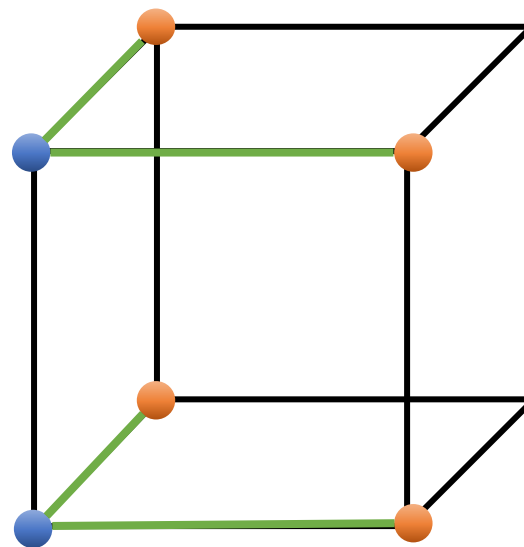
And

$$I_f = \sum_{i=1}^n \sum_{S \ni i} \hat{f}(S)^2 = \sum_S |S| \hat{f}(S)^2$$

# Isoperimetric Inequalities for the Hypercube

Let  $G = (V, E)$  be an undirected graph. For  $A \subseteq V$ , the **edge boundary** of  $A$  is the set  $E(A, A^c)$  of edges with one endpoint in  $A$  and one in  $V \setminus A$ . The **vertex boundary** of  $A$  is the set

$$\partial A = \{v \in A^c : v \text{ has a neighbor in } A\}$$



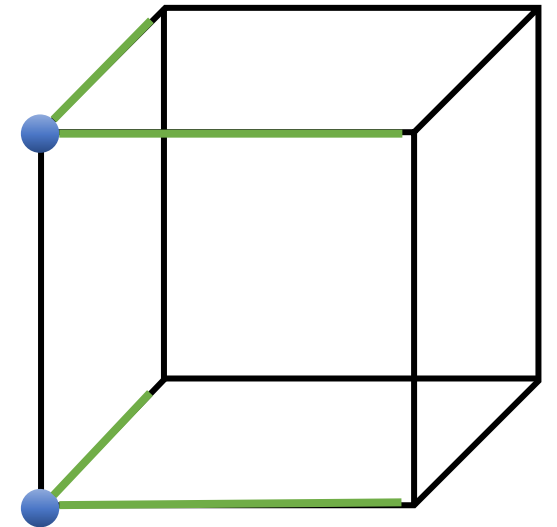
# Isoperimetric Inequalities for the Hypercube

**Theorem (Harper's edge isoperimetric inequality).** Every subset  $A$  of the vertices of the hypercube  $Q_n$  satisfies

$$|E(A, A^c)| \geq |A| \log \frac{2^n}{|A|}$$

with equality when  $A$  is a subcube.

- $n = 3$ 
  - $|A| = 4, |E(A, A^c)| \geq 4 * \log \frac{8}{4} = 4$
  - $|A| = 2, |E(A, A^c)| \geq 2 * \log \frac{8}{2} = 4$



# Isoperimetric Inequalities for the Hypercube

**Theorem (Harper's vertex isoperimetric inequality).** Every subset  $A$  of the vertices of the hypercube  $Q_n$  with  $\binom{n}{\leq r} \leq |A| < \binom{n}{\leq r+1}$  satisfies

$$|A \cup \partial A| \geq \binom{n}{\leq r+1}$$

where  $\binom{n}{\leq r} := \sum_{i=0}^r \binom{n}{i} = |B_r(x)|$  (the size of a Hamming ball of radius  $r$ )

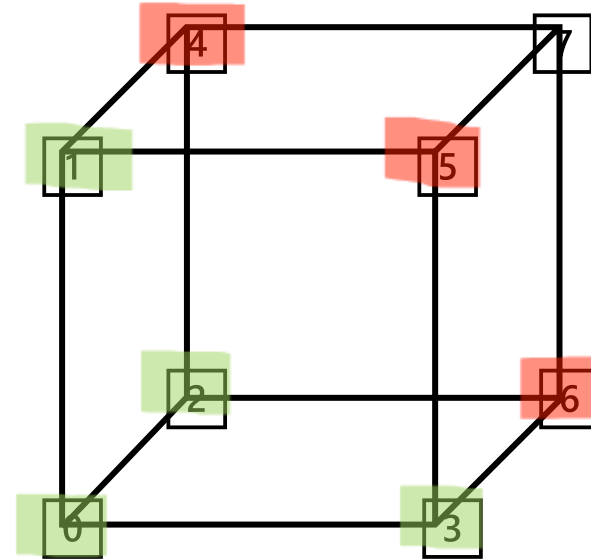
Define the **simplicial order** of the hypercube:

$$000 < 100 < 010 < 001 < 110 < 101 < 011 < 111$$

**Theorem (Harper's vertex isoperimetric inequality).** For any  $A$ ,

$$|\partial A| \geq |\partial L_{|A|}|,$$

where  $L_{|A|}$  is the initial segment in simplicial orderings of length  $|A|$



# Influence and Isoperimetric Inequalities for the Hypercube

For a Boolean function  $f: \{0,1\}^n \rightarrow \{0,1\}$ , we can define:

- **Vertex boundary:**  $V_f := \{x \in \{0,1\}^n : s_f(x) > 0\}$
- **Edge boundary:**  $E_f := \{(x, x + e_i) : f(x) \neq f(x + e_i)\}$

**Claim.**  $I_f = |E_f|/2^{n+1}$

*Proof.*

- Notice that  $2|E_f| = \sum_x s_f(x)$
- Thus,

$$I_f = \frac{1}{4} \mathbb{E}[s_f(x)] = 2^{-(n+2)} \cdot 2|E_f| = 2^{-(n+1)} |E_f|$$



# Influence and Isoperimetric Inequalities for the Hypercube

**Theorem (Harper's edge isoperimetric inequality).** Every  $f: \{0,1\}^n \rightarrow \{0,1\}$  satisfies

$$I_f \geq \frac{1}{2} \mathbb{E}[f] \log \frac{1}{\mathbb{E}[f]}$$

*Proof.*

- $f: \{0,1\}^n \rightarrow \{0,1\} \iff f = \mathbf{1}_A$  where  $A := \{x : f(x) = 1\}$
- $|E_f| = |E(A, A^c)| \geq |A| \log \frac{2^n}{|A|} = 2^n \mathbb{E}[f] \log(\mathbb{E}[f]^{-1})$
- $I_f = |E_f| \cdot 2^{-(n+1)}$



# Influence and Isoperimetric Inequalities for the Hypercube

Poincare inequality:

$$\frac{|E_f|}{2^{n+1}} = I_f \geq \text{Var}[f]$$

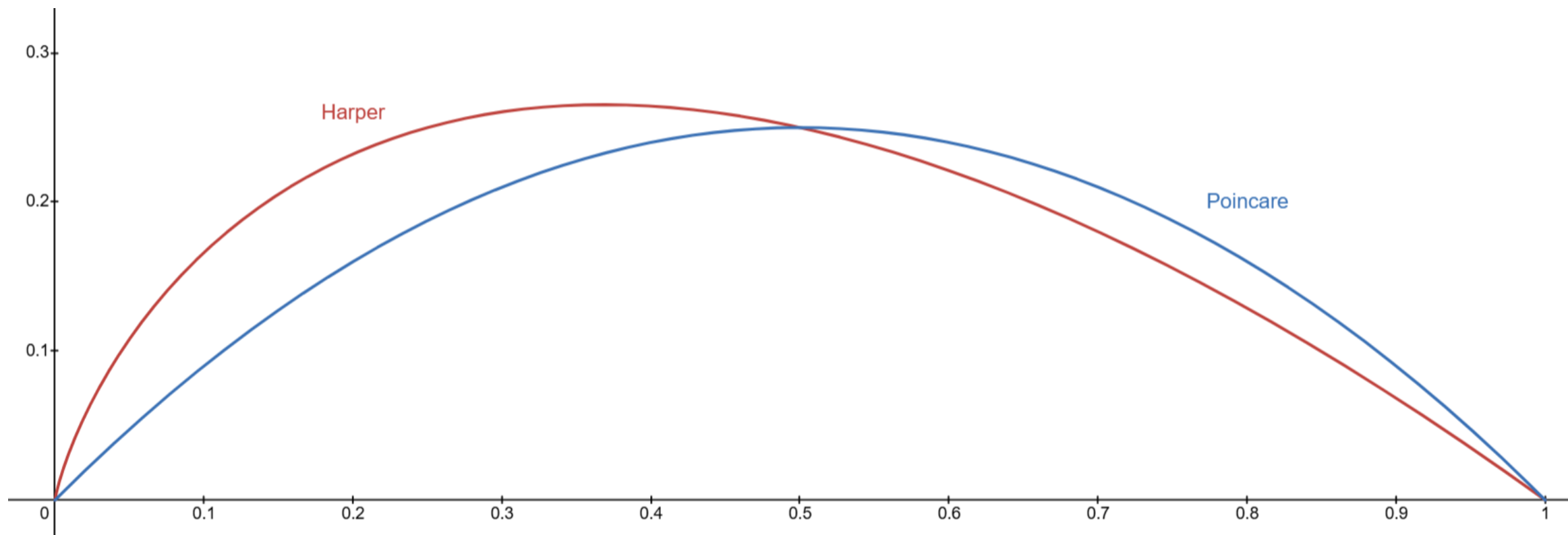
*Proof.*

- $I_f = \sum_S |S| \hat{f}(S)^2 \geq \sum_{S \neq \emptyset} \hat{f}(S)^2 = \text{Var}[f]$



# Influence and Isoperimetric Inequalities for the Hypercube

$$I_f \geq \max \left\{ \text{Var}[f], \frac{1}{2} \mathbb{E}[f] \log \frac{1}{\mathbb{E}[f]} \right\}$$



# Influence and Isoperimetric Inequalities for the Hypercube

What about the **vertex boundary**  $V_f := \{x \in \{0,1\}^n : s_f(x) > 0\}$ ?

## Examples:

- Dictator  $f(x) = x_1$ :  $V_f = \{0,1\}^n$ ,  $|V_f| = 2^n$
- Majority  $f(x) = \text{Majority}(x)$ :

$$V_f = \left\{x \in \{0,1\}^n : |x|_1 = \frac{n}{2}\right\}, \quad |V_f| = \binom{n}{n/2} = \mathcal{O}\left(\frac{2^n}{\sqrt{n}}\right)$$

This is optimal by the Kruskal-Katona theorem

**Theorem (Margulis '74).** For every  $f: \{0,1\}^n \rightarrow \{0,1\}$ ,

$$\frac{|V_f|}{2^n} \cdot \frac{|E_f|}{2^n} \simeq \frac{|V_f|}{2^n} \cdot I_f \geq \Omega(\text{Var}[f]^2)$$

- When the vertex boundary is small, edge boundary (or influence) must be very large

# Influence and Isoperimetric Inequalities for the Hypercube

Michel Talagrand considered a new quantity  $\mathbb{E}[\sqrt{s_f(x)}]$  (Talagrand boundary of  $f$ ) and proved several results

- The basic one is  $\mathbb{E}[s_f(x)^{1/2}] \geq \Omega(\text{Var}[f])$
- By Cauchy-Schwarz inequality,

$$\mathbb{E}[s_f(x)^{1/2}] = \mathbb{E}[s_f(x)^{1/2} \cdot \mathbf{1}_{s_f(x)>0}] \leq \mathbb{E}[s_f(x)]^{1/2} \cdot \Pr[s_f(x) > 0]^2 \simeq \left( \frac{|V_f|}{2^n} \cdot I_f \right)^{1/2}$$

recovering Margulis's theorem

**Theorem (Talagrand '93).** For all non-constant  $f: \{0,1\}^n \rightarrow \{0,1\}$ ,

$$\mathbb{E}[s_f(x)^{1/2}] \geq \Omega\left(\text{Var}[f] \log \frac{1}{\text{Var}[f]}\right)$$

- Tight for Majority function

# Today's Lecture

- Influence and Isoperimetric Inequalities for the Hypercube
- **KKL Theorem**
- Unifying Talagrand and KKL

# KKL Theorem

**Theorem (Kahn-Kalai-Linial '88).** For  $f: \{0,1\}^n \rightarrow \{0,1\}$ , let  $I_{\max} := \max_i I_i(f)$ . Then,

$$\text{Var}[f] \lesssim \frac{I_f}{\log(1/I_{\max})}$$

Since  $I_f \leq nI_{\max}$ , we have

$$I_{\max} \geq \Omega\left(\frac{\log n}{n} \text{Var}[f]\right)$$

- KKL is another strengthening of Poincare:

$$I_f \gtrsim \min_i \log\left(\frac{1}{I_i(f)}\right) \cdot \text{Var}[f]$$

- If all individual influences of  $f$  being small, then it gains a large factor

# KKL Theorem

**Theorem (Kahn-Kalai-Linial '88).** For  $f: \{0,1\}^n \rightarrow \{0,1\}$ , let  $I_{\max} := \max_i I_i(f)$ . Then,

$$\text{Var}[f] \lesssim \frac{I_f}{\log(1/I_{\max})}$$

*Proof.*

- Recall that  $\text{Var}[f] = \sum_{S \neq \emptyset} \hat{f}(S)^2$  and  $I_f = \sum_S |S| \hat{f}(S)^2$
- Bounding high-degree coefficients:

$$I_f \geq k \cdot \sum_{|S| > k} \hat{f}(S)^2, \quad \Rightarrow \quad \sum_{|S| > k} \hat{f}(S)^2 \leq \frac{I_f}{k}$$

- Bounding low-degree coefficients:

$$\sum_{|S| \leq k} |\hat{f}(S)|^2 \leq \sum_{i \in [n]} \sum_{i \in S, |S| \leq k} |\hat{f}(S)|^2 = \sum_{i \in [n]} \|(\partial_i f)^{\leq k}\|_2^2$$

where  $\partial_i f(x) = \sum_{i \in S} \hat{f}(S) \chi_S(x)$  and  $(\partial_i f)^{\leq k}$  takes the degree  $\leq k$  terms

# Detour: Hypercontractivity

**Theorem.** For  $f: \{0,1\}^n \rightarrow \mathbb{R}$  and  $k > 0$ ,

1) For  $2 \leq q < \infty$ ,

$$\|f^{\leq k}\|_q \leq (q-1)^{\frac{k}{2}} \|f\|_2$$

2) For  $1 \leq p \leq 2$ ,

$$\|f^{\leq k}\|_2 \leq e^{\left(\frac{2}{p}-1\right)k} \|f\|_p$$

• 2) can be simplified: let  $q$  be such that  $\frac{1}{q} + \frac{1}{p} = 1$ . Then by Hölder's and 1),

$$\|f^{\leq k}\|_2^2 = \langle f^{\leq k}, f \rangle \leq \|f^{\leq k}\|_q \|f\|_p \leq (q-1)^{\frac{k}{2}} \|f^{\leq k}\|_2 \|f\|_p = (p-1)^{-\frac{k}{2}} \|f^{\leq k}\|_2 \|f\|_p$$

$$\|f^{\leq k}\|_2 \leq (p-1)^{-\frac{k}{2}} \|f\|_p$$

# KKL Theorem

**Theorem (Kahn-Kalai-Linial '88).** For  $f: \{0,1\}^n \rightarrow \{0,1\}$ , let  $I_{\max} := \max_i I_i(f)$ . Then,

$$\text{Var}[f] \lesssim \frac{I_f}{\log(1/I_{\max})}$$

*Proof.*

- **Bounding low-degree coefficients:**

$$\sum_{|S| \leq k} |\hat{f}(S)|^2 \leq \sum_{i \in [n]} \sum_{i \in S, |S| \leq k} |\hat{f}(S)|^2 = \sum_{i \in [n]} \|(\partial_i f)^{\leq k}\|_2^2$$

- By hypercontractivity,

$$\sum_{|S| \leq k} |\hat{f}(S)|^2 \leq \sum_{i \in [n]} \|(\partial_i f)^{\leq k}\|_2^2 \leq \sum_{i \in [n]} 3^k \|\partial_i f\|_{4/3}^2$$

- **Claim:**  $\mathbb{E}[|\partial_i f|^p] \leq \mathbb{E}[|\partial_i f|] = 2\mathbb{E}[|\partial_i f|^2] = 2I_i(f)$

- $\sum_{i \in [n]} \|\partial_i f\|_{4/3}^2 \simeq \sum_{i \in [n]} I_i(f)^{3/2} \leq I_{\max}^{1/2} \cdot \sum_{i \in [n]} I_i(f) = I_{\max}^{1/2} \cdot I_f$

# KKL Theorem

**Theorem (Kahn-Kalai-Linial '88).** For  $f: \{0,1\}^n \rightarrow \{0,1\}$ , let  $I_{\max} := \max_i I_i(f)$ . Then,

$$\text{Var}[f] \lesssim \frac{I_f}{\log(1/I_{\max})}$$

*Proof.*

- Thus, we have

$$\text{Var}[f] = \sum_{|S| > k} \hat{f}(S)^2 + \sum_{|S| \leq k} |\hat{f}(S)|^2 \lesssim \frac{I_f}{k} + 3^k I_{\max}^{1/2} \cdot I_f$$

- Taking  $k \simeq \log \frac{1}{I_{\max}}$  proves the theorem



# KKL Theorem

**Theorem (Kahn-Kalai-Linial '88).** For  $f: \{0,1\}^n \rightarrow \{0,1\}$ , let  $I_{\max} := \max_i I_i(f)$ . Then,

$$I_{\max} \geq \Omega\left(\frac{\log n}{n} \text{Var}[f]\right)$$

KKL is tight due to the **tribes function**:

$$f(x) = \bigvee_{i \in [m]} \bigwedge_{j \in [k]} x_{i,j}, \quad k = \log n - \log \log n, m = n/k$$

- $\text{Var}[f] \geq \Omega(1)$
- $I_i(f) \leq (1 - o(1)) \frac{\ln n}{2n}$

# Today's Lecture

- Influence and Isoperimetric Inequalities for the Hypercube
- KKL Theorem
- **Unifying Talagrand and KKL**

# Unifying Talagrand and KKL

$$\mathbb{E}[s_f(x)^{1/2}] \geq \Omega\left(\text{Var}[f] \log \frac{1}{\text{Var}[f]}\right)$$

$$I_{\max} \geq \Omega\left(\frac{\log n}{n} \text{Var}[f]\right)$$

Talagrand showed that there **exists** an  $\alpha \in (0, 1/2]$  such that

$$\mathbb{E}[s_f(x)^{1/2}] \gtrsim \text{Var}[f] \cdot \log^{1/2-\alpha}\left(\frac{1}{\text{Var}[f]}\right) \cdot \log^\alpha\left(1 + \frac{1}{\sum_i I_i(f)^2}\right)$$

- $\mathbb{E}[s_f(x)^{1/2}] \leq \mathbb{E}[s_f(x)]^{1/2} = I_f^{1/2}$

- $\frac{1}{\sum_i I_i(f)^2} \geq \frac{1}{I_{\max} I_f}$

- It recovers KKL if  $\alpha = 1/2$



$$I_{\max} \geq e^{-o\left(\left(\frac{I_f}{\text{Var}[f]}\right)^{\frac{1}{2\alpha}}\right)}$$

# Unifying Talagrand and KKL

$$\mathbb{E}[s_f(x)^{1/2}] \geq \Omega\left(\text{Var}[f] \log \frac{1}{\text{Var}[f]}\right)$$

$$I_{\max} \geq \Omega\left(\frac{\log n}{n} \text{Var}[f]\right)$$

Talagrand showed that there **exists** an  $\alpha \in (0, 1/2]$  such that

$$\mathbb{E}[s_f(x)^{1/2}] \gtrsim \text{Var}[f] \cdot \log^{\frac{1}{2}-\alpha}\left(\frac{1}{\text{Var}[f]}\right) \cdot \log^\alpha\left(1 + \frac{1}{\sum_i I_i(f)^2}\right)$$

**Theorem (Eldan-Gross '20).**

$$\mathbb{E}[s_f(x)^{1/2}] \gtrsim \text{Var}[f] \cdot \sqrt{\log\left(1 + \frac{1}{\sum_i I_i(f)^2}\right)}$$

# Unified Proof by Eldan-Kindler-Lifshitz-Minzer

If we define  $\nabla f: \{0,1\}^n \rightarrow \{-1,0,1\}^n$  to be  $\nabla f(x) = (\partial_1 f(x), \partial_2 f(x), \dots, \partial_n f(x))$ , then

$$\|\nabla f(x)\|_2 = \sqrt{s_f(x)}$$

Since  $|\partial_i f| = |f(x_{i=0}) - f(x_{i=1})|/2 = |\sum_{i \in S} \hat{f}(S) \chi_{S-\{i\}}|$ , we have

$$\mathbb{E}[|\partial_i f|] = \mathbb{E}\left[\left|\sum_{i \in S} \hat{f}(S) \chi_{S-\{i\}}\right|\right] \geq \left|\sum_{i \in S} \hat{f}(S) \mathbb{E}[\chi_{S-\{i\}}]\right| = |\hat{f}(\{i\})|$$

Thus,

$$\mathbb{E}[\|\nabla f(x)\|_2] \geq \left\|(\mathbb{E}[|\partial_i f|])_{i \in [n]}\right\|_2 \geq \left\|(|\hat{f}(\{i\})|)_{i \in [n]}\right\|_2 = \left(\sum_i \hat{f}(\{i\})^2\right)^{1/2} = \sqrt{\mathbf{w}^1[f]}$$

That is,

$$\mathbb{E}[s_f(x)^{1/2}] \geq \sqrt{\mathbf{w}^1[f]} \geq \mathbf{w}^1[f]$$

# Unified Proof

Suppose that we knew  $f$  has significant weight on level- $d$  ( $d > 1$ ); can we use the same logic to lower bound  $\mathbb{E}[\|\nabla f\|_2]$ ?

## Random Restriction

If  $f$  has sizable mass around level  $d$ , then  $f_{\bar{J} \rightarrow z}$  has sizable weight around level  $pd$

**Corollary.** Let  $d \in \mathbb{N}$  and  $p \in [0,1]$  be such that  $pd \geq 10$ . Suppose that  $f: \{0,1\}^n \rightarrow \{-1,1\}$ , and let  $(J, z)$  be a  $p$ -random restriction. Define the weight around level  $d$  to be

$$\mathbf{w}^{\approx d}[f] := \sum_{d \leq k < 2d} \mathbf{w}^k[f]$$

Then

$$\mathbb{E}_{J,z} [\mathbf{w}^{\approx pd}[f_{\bar{J} \rightarrow z}]] \geq \Omega(\mathbf{w}^{\approx d}[f])$$

- If we take  $p = \frac{1}{2d}$ , then  $\mathbb{E}_{J,z} [\mathbf{w}^1[f_{\bar{J} \rightarrow z}]] \gtrsim \mathbf{w}^{\approx d}[f]$

# Unified Proof

**Lemma.** Let  $d \in \mathbb{N}$ , and  $f: \{0,1\}^n \rightarrow \{0,1\}$ . Then  $\mathbb{E}[\|\nabla f\|_2] \gtrsim \sqrt{d} \mathbf{W}^{\approx d}[f]$

*Proof.*

- Let  $(J, z)$  be a  $1/2d$ -random restriction. Then, for any  $x \in \{0,1\}^n$ , we have

$$\mathbb{E}_J \left[ \|\nabla f_{\bar{J} \rightarrow z}(x_J)\|_2 \right] \leq \sqrt{\mathbb{E}_J \left[ \|\nabla f_{\bar{J} \rightarrow z}(x_J)\|_2^2 \right]} = \sqrt{\mathbb{E}_J \left[ \sum_j |\partial_j f(x)|^2 \cdot 1_{j \in J} \right]} = \frac{\|\nabla f\|_2}{\sqrt{2d}}$$

Thus,

$$\mathbb{E}[\|\nabla f\|_2] \gtrsim \sqrt{d} \mathbb{E}_{J,z} \left[ \mathbb{E}_x \left[ \|\nabla f_{\bar{J} \rightarrow z}(x_J)\|_2 \right] \right] \geq \sqrt{d} \mathbb{E}_{J,z} \left[ \mathbf{W}^1[f_{\bar{J} \rightarrow z}] \right] \gtrsim \sqrt{d} \mathbf{W}^{\approx d}[f]$$



- If we apply the lemma for  $d, 2d, 4d, \dots$  and sum them up, we can further bootstrap:

$$\mathbb{E}[\|\nabla f\|_2] \gtrsim \sqrt{d} \mathbf{W}^{\geq d}[f]$$

# Unified Proof

Then, we can apply the known **Fourier tail bounds** to prove Talagrand and Eldan-Gross:

- If  $f$  has **small variance**, then most of its Fourier mass lies above level  $d = \Omega(\log(1/\text{Var}[f]))$
- If  $f$  has **small influences**, then most Fourier mass lies above level  $d = \Omega(\log(1/\sum_i I_i(f)^2))$

# Unified Proof

$$\mathbb{E}[s_f(x)^{1/2}] \geq \Omega\left(\text{Var}[f] \log \frac{1}{\text{Var}[f]}\right)$$

*Proof.*

Case 1:  $\text{Var}[f] \geq 2^{-16}$ :

$$\mathbb{E}[\|\nabla f\|_2] \gtrsim \mathbf{W}^{\geq 1}[f] = \text{Var}[f]$$

Case 2:  $\text{Var}[f] < 2^{-16}$ :

$$\sum_{0 < |S| \leq d} \hat{f}(S)^2 \leq \|f^{\leq d}\|_2^2 = \langle f^{\leq d}, f \rangle \leq \|f^{\leq d}\|_4 \|f\|_{4/3}$$

• Hypercontractivity:  $\|f^{\leq d}\|_4 \leq 3^{d/2} \|f^{\leq d}\|_2 \leq 3^{d/2} \|f\|_2$

$$\|f^{\leq d}\|_4 \|f\|_{4/3} \leq 3^{d/2} \|f\|_2 \|f\|_{4/3} = 3^{d/2} \mathbb{E}[f]^{5/4}$$

• We may assume  $\mathbb{E}[f] \leq 1/2$ , and thus,  $\mathbb{E}[f] \leq 2\text{Var}[f]$

• If we take  $d = \frac{1}{8} \log(1/\text{Var}[f])$ , then  $\sum_{0 < |S| \leq d} \hat{f}(S)^2 \leq 0.9\text{Var}[f]$  and  $\mathbf{W}^{\geq d}[f] \geq 0.1\text{Var}[f]$

# Unified Proof

$$\mathbb{E}[s_f(x)^{1/2}] \gtrsim \text{Var}[f] \cdot \sqrt{\log\left(1 + \frac{1}{\sum_i I_i(f)^2}\right)}$$

*Proof.*

**Theorem (Keller-Kindle '13).** There are  $c_1, c_2 > 0$  such that for any  $f: \{0,1\}^n \rightarrow \{0,1\}$ ,

$$\mathbf{W}^{\leq c_1 \log(1/M[f])}[f] \leq M[f]^{c_2}$$

where  $M[f] := \sum_i I_i(f)^2$

Case 1:  $M[f] \geq \text{Var}[f]^{2/c_2}$

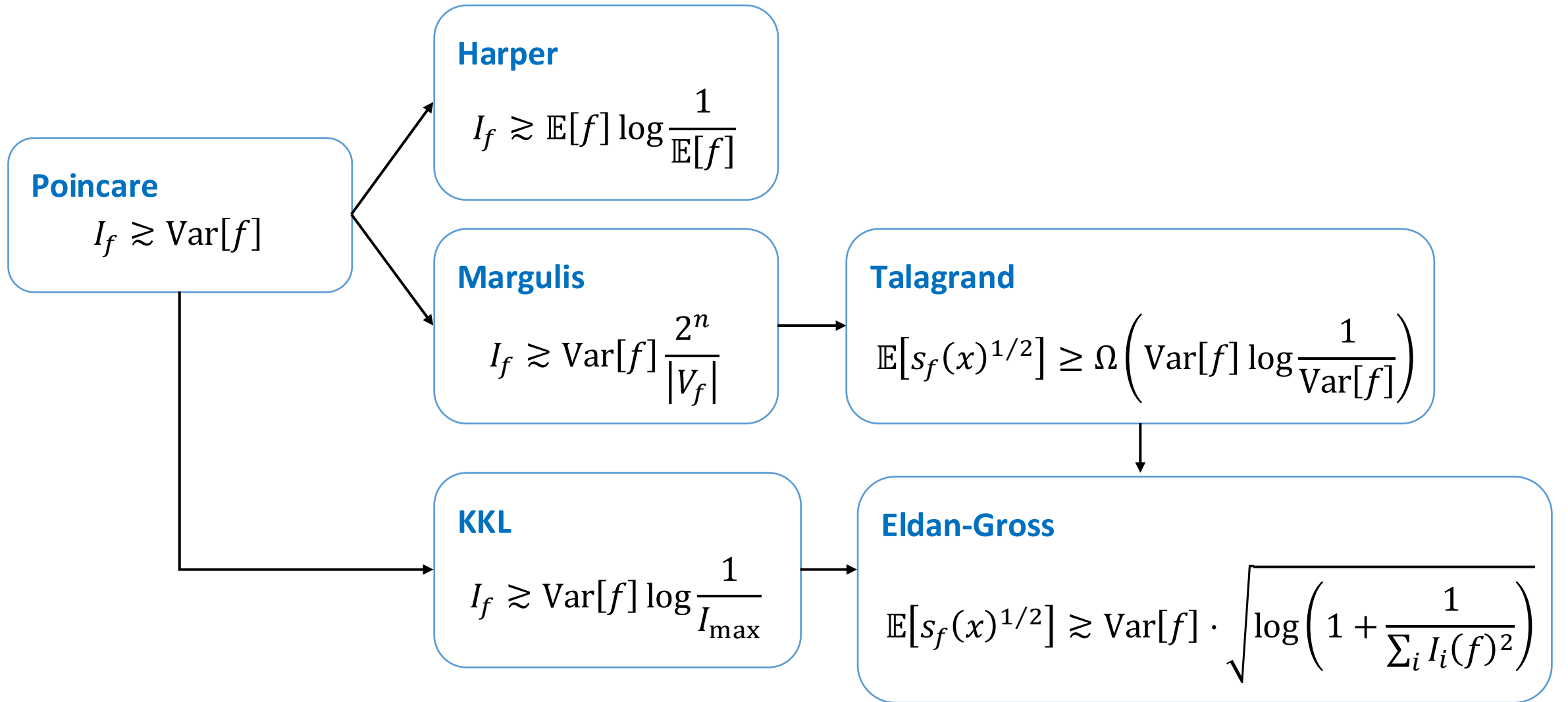
$$\text{Var}[f] \cdot \log^{1/2}\left(1 + \frac{1}{M[f]}\right) \lesssim \text{Var}[f] \log \frac{1}{\text{Var}[f]}$$

Case 2:  $M[f] < \text{Var}[f]^{2/c_2}$

- Take  $d = c_1 \log(1/M[f])$ , and we have

$$\mathbf{W}^{\leq d}[f] \leq M[f]^{c_2} \leq \text{Var}[f]^2 \leq 0.5\text{Var}[f]$$





# Final Story: Aaronson-Ambainis Conjecture

**Conjecture.** Any function  $f: \{-1, 1\}^n \rightarrow [-1, 1]$  of degree  $d$  satisfies

$$I_{\max} \geq \text{poly} \left( \frac{\text{Var}[f]}{d} \right)$$

- If any quantum algorithm has quantum query complexity  $q$ , then there exists a classical algorithm that makes  $\text{poly}(q)$  queries and succeeds on average over the inputs
- **KKL**  $\implies$  for **Boolean**  $f$ ,  $I_{\max} \gtrsim \frac{\log n}{n} \text{Var}[f]$ .  
Since  $D(f) \leq \text{deg}(f)^4$  and  $f$  can only depend on at most  $2^{D(f)}$  many variables, it gives an **exponential bound** for AA conjecture
- **(O'Donnell-Saks-Schrammand-Servedio '05)** proves that  $I_{\max} \gtrsim \frac{\text{Var}[f]}{D(f)} \geq \frac{\text{Var}[f]}{d^4}$ , which implies a **polynomial bound** for AA conjecture for Boolean functions
- Read Chap. 10 of Thomas Rothvoss's lecture notes for more discussion